

# Discrete Mathematics

Second Stage  
Information Technology Department  
Lecture 2

# Introduction

- Mathematics is divided into two major branches discrete mathematics and continuous mathematics.
- Discrete mathematics deals with only those real numbers which are multiples of same basic unit. If the basic unit is 1, then the discrete variable can assume only integral values.
- Thus in continuous mathematics, a number system is usually real numbers while for discrete mathematics it is the integer.

# Introduction (Cont'd...)

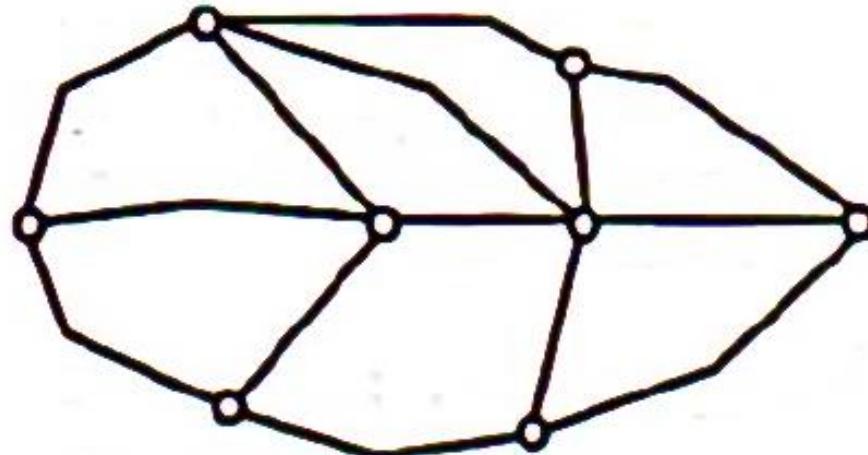
- Discrete Mathematics is mathematics that deals with discrete objects.
- Discrete objects are those which are separated from (not connected to/distinct from) each other.

# Introduction (Cont'd...)

- Discrete mathematics does not support the notion of continuity.
- Discrete mathematics discusses languages used in mathematical reasoning, basic concepts, and their properties and relationships.

# Introduction (Cont'd...)

- Examples of discrete mathematics:
  - Determine the **minimum number** of colors needed to **color** the **regions** in the **diagram** given below and explain why this number is **minimum**.



# Propositional Calculus

# Statements (or Propositions)

- Statements are kinds of sentences that we have to use to convey our thoughts to others.
- A **sentence** is a statement or proposition if it is possible to say whether what is conveyed by the sentence is true or false.
- Statements are **logical** entities.
- Sentences are **grammatical** entities.
- Not all sentences **express** statements and some sentences may **express** more than one statement.

# Statements (or Propositions)

## Examples

- Mars is a planet.
- $9 > 13$ .
- $Y + 8 = 12$
- There are 12 months in a year.

# Statements (or Propositions) cont'd...

- A **statement** is a declarative sentence that is either **true** or **false**.
- Hence examples **(1)** and **(4)** *are statements*.
- In Example **(3)** ,the **statement** is true depends on the value of **y**.
- If **Y** is **4** the sentence is true, if  $y \neq 4$  then the sentence is false.

# Open Statement

- In example (3) , if we put  $y = 4$ , it becomes a true statement, if we take value of  $y \neq 4$ , it becomes false. Such statements are **open statements**.
- Thus if a mathematical sentence is neither true nor false it is called **open sentence**.
- An **open statement** is a sentence that contains one or more **variables** such that when certain **values** are substituted for the variables, we get statements.

# Truth Table

- A table giving all possible truth values of a statement is called **truth table**.
- Statement has a definite **truth value** which is either **true or false**.
- True values are denoted by (T) and false values are denoted by (P).

# Exercise

1. *Which of the following statements are true and which are false ?*
  - a)  $9 < 12$
  - b)  $2 + 5 = 3 + 9$
2. *What type of this sentence is*
  - $x + 8 = 17 ?$

## Exercise (Cont'd...)

3. *For what value of  $x$  following sentences will become true statements ?*

- a)  $3x+9=15$
- b)  $x+6=8$
- c)  $x+1>5$
- d)  $x+2<8$
- e)  $5x \geq 25$
- f)  $5x \leq 25$

## Exercise (Cont'd...)

4. *Which of the following statements are true?*

- a)  $x+4=6$       when  $x=2$
- b)  $x+4 \neq 6$       when  $x=2$
- c)  $x + 5 \neq 8$       when  $x = 3$
- d)  $2x + 4y = 14$       when  $x = 1, Y = 3$
- e)  $3x + 5y = 11$       when  $x = 0, y = 2$

# Answers

# Logical Connectives

- Every statement must be either true or false but not both.
- If two or more statements, they can be combined to produce a new statement.
- These new statements are called compound statements.
- To combine statements we use difference symbols.

# Conjunction ('^' or 'and')

- If  $p$  and  $q$  are statements, the compound statement ' $p \wedge q$ ' is ' $p$  and  $q$ ' and is called ' $p$  conjunction  $q$ ' or ' $p$  meet  $q$ '.
- ' $\wedge$ ' denotes 'and' and is known as **conjunction**.
- **Examples**
- Let us consider the statements.
  - $p$ :  $3 > 2$
  - $q$  :  $9$  is an odd number,
  - then  $p \wedge q$  is the statement
  - $3 > 2$  and  $9$  is an odd number.

# Truth table for ' $p \wedge q$ '

<b>p</b>	<b>q</b>	<b><math>p \wedge q</math></b>
T	T	T
F	F	F
T	F	F
F	T	F

- ' $p \wedge q$ ' is true only when both are true otherwise false.

# Disjunction (V or 'or')

- When **two or more statements** are combined by the word **'or'**, the compound statement is known as **'disjunction'**.
- The symbol ' $p \vee q$ ' is read as ' $p$  or  $q$ ' or ' $p$  disjunction  $q$ ' or ' $p$  join  $q$ '.
- Example:
  - $p$ :  $a$  is equal to 5.
  - $q$  :  $b$  is equal to 7,
  - then  $P \vee q$  is the statement
  - $a$  is equal to 5 or  $b$  is equal to 7.

# Truth table for $p \vee q$

$p$	$q$	$p \vee q$
T	T	T
T	F	T
F	T	T
F	F	F

- ' $p \vee q$ ' is true if either  $p$  is true or  $q$  is true or both  $p$  and  $q$  are true and  $p \vee q$  is false if both  $p$  and  $q$  are false.

# Remark

- The exclusive disjunction or exclusive 'OR' of two propositions  $p$  and  $q$  is the statement. Either  $p$  is true or  $q$  is true, but both are not true. Either  $p$  is true or  $q$  is true, but both are not true, we denote this by  $p \oplus q$ .

# Truth table for $p \oplus q$ (exclusive 'OR')

$p$	$q$	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

# Negation ( $\sim P$ )

- Let  $P$  be any statement then negation of  $p$  is denoted by ' $\sim P$ ' (or  $\bar{P}$  ) is read as 'not  $P$ '
- If  $P$  is true then  $\sim P$  is false.
- If  $p$  is false then  $\sim P$  is true,
- Truth table**

$P$	$\sim P$
T	F
F	T

- $\sim P$  is a unary connective as only one statement is required to form negation.

# Example

## Example (1):

- If  $P$  is the statement 'Aram is intelligent boy',
- then  $\sim p$  is the statement 'Aram is not intelligent boy.'

## Example (2):

- If  $P$  is the statement 'I like to read'.
- then  $\sim p$  is the statement 'I don't like to read'

# Homework

- Find truth table of the following proposition:
  1.  $(\sim p \wedge (\sim q \wedge p)) \vee (q \wedge p) \vee (p \wedge p)$
  2.  $(p \wedge q) \vee (\sim p \wedge q) \vee (p \wedge \sim q) \vee (\sim p \wedge \sim q)$

# Conditional Proposition

- If ***p*** and ***q*** are proposition. the compound proposition "if *p* then *q*" denoted by  $p = > q$  is called a **conditional proposition**
- Example:
- If it rains then I will carry an umbrella.
- Here:
  - $p$  : *It rains* ( antecedent )
  - $q$  : *I will carry an umbrella* (consequent)

- The connective **if .. .... ... then** can also be read as follows.
  1.  $p$  *Implies*  $q$ .
  2.  $p$  is *sufficient* for  $q$ .
  3.  $p$  *only if*  $q$ .
  4.  $q$  is *necessary* for  $p$ .
  5.  $q$  *if*  $p$ .
  6.  $q$  *follows from*  $p$ .
  7.  $q$  is *consequence* of  $p$ .

- The truth table for implication is given in Table:

$p$	$q$	$p \Rightarrow q$
T	T	T
T	F	F
F	T	T
F	F	T

- Note : The only circumstances under which the implication  $p \Rightarrow q$  is *false* when  $p$  is *true* and  $q$  is *false*.

- Example: Calculate truth table for:

$$(i) \ p \vee \sim q \Rightarrow p$$

$$(ii) ((\sim (p \wedge q) \vee r) \Rightarrow \sim p)$$

- I. Solution: The truth of the given compound statement is shown below.

$p$	$q$	$\sim q$	$p \vee \sim q$	$p \vee \sim q \Rightarrow p$
T	T	F	T	T
T	F	T	T	T
F	T	F	F	T
F	F	T	T	F

- II. Homework.....

# Converse, Contrapositive and Inverse

- There are some related implication that can be formed from  $p \Rightarrow q$ . *If  $p \Rightarrow q$  is an implication.*
- Then the **converse** of  $p \Rightarrow q$  is the implication  $q \Rightarrow p$ .
- *The contrapositive of  $p \Rightarrow q$  is the implication  $\sim q \Rightarrow \sim p$  .*
- *The inverse of  $p \Rightarrow q$  is  $\sim p \Rightarrow \sim q$ .*

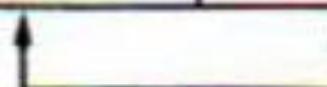
- The truth table of the four propositions follow:

		Conditional $p \Rightarrow q$	Converse $q \Rightarrow p$	Inverse $\neg p \Rightarrow \neg q$	Contrapositive $\neg q \Rightarrow \neg p$
$p$	$q$				
T	T	T	T	T	T
T	F	F	T	T	F
F	T	T	F	F	T
F	F	T	T	T	T

- Example:
- Consider the statement
  - $p$  : *It rains.*
  - $q$  : *The crops will grow*
- The implication ( $p \Rightarrow q$ ) states that.
  - If it rains then the crops will grow.
- The converse ( $q \Rightarrow p$ ) states that.
  - *If the crops grow. then there has been rain.*
- The contrapositive  $\sim q \Rightarrow \sim p$  states that.
  - *if the crops do not grow then there has been no rain.*
- The inverse  $\sim p \Rightarrow \sim q$  states that.
  - *If it does not rain then the crops will not grow.*

- Example : Show that contrapositives are logically equivalent; that is :  $\neg q \Rightarrow \neg p \equiv p \Rightarrow q$
- Solution:

$p$	$q$	$\neg p$	$\neg q$	$\neg q \Rightarrow \neg p$	$p \Rightarrow q$
T	T	F	F	T	T
T	F	F	T	F	F
F	T	T	F	T	T
F	F	T	T	T	T



- Homework: Prove that if  $x^2$  is divisible by 4, then  $x$  is even.

# Biconditional Statement

- If  $p$  and  $q$  are statement, then the compound statement  $p$  if and only if  $q$ . denoted by  $p \leftrightarrow q$  is called a biconditional statement.
- The truth table of  $p \leftrightarrow q$  is given:

$p$	$q$	$p \leftrightarrow q$
T	T	T
T	F	F
F	T	F
F	F	T

- Example : Show that  $p \Leftrightarrow q \equiv (p \Rightarrow q) \wedge (q \Rightarrow p)$

Solution:

$p$	$q$	$p \Leftrightarrow q$	$p \Rightarrow q$	$q \Rightarrow p$	$(p \Rightarrow q) \wedge (q \Rightarrow p)$
T	T	T	T	T	T
T	F	F	F	T	F
F	T	F	T	F	F
F	F	T	T	T	T



- Homework: Show that  $p \Leftrightarrow q \equiv (p \vee q) \Rightarrow (p \wedge q)$

# Negation of Compound Statement

- Negation of Conjunction : the negation of the conjunction, would mean the negation of at least one of the two sub-statements.
- $\sim(p \wedge q) \equiv \sim p \vee \sim q$  .
- Example : Write the negation of each of the following conjunctions: ( $2 + 4 = 6$  and  $7 < 12$  ).
- Solution: Let  $p : 2 + 4 = 6$  and  $q : 7 < 12$ .
- Then the conjunction is given by " $p \wedge q$ ".
- Now :  $\sim p : 2 + 4 \neq 6$  and  $\sim q : 7 \geq 12$ .
- The negation of " $p \wedge q$ " is given by
- $\sim(p \wedge q) : 2 + 4 \neq 6$  or  $7 \geq 12$

# Negation of Disjunction

- The **negation** of a **disjunction**  $p \vee q$  is the **conjunction of the negation of p and the negation of q**. we write:  $\sim(p \vee q) = \sim p \wedge \sim q$
- **Example :** Write the negation of each of the following disjunction: 9 is greater than 4 or 6 is less than 8.
- **Solution:** Let  $p$  : 9 is greater than 4  
 $q$  : 6 is less than 8.
- Then negation of  $p \vee q$  is given by
- $\sim(p \vee q)$  9 is not greater than 4 and 6 is not less than 8.

# Negation of a Negation

- A negation of negation of a statement is the statement itself. Equivalently. we write  $\sim(\sim p) \equiv p$ .
- **Example :** Verify for the statement (*Roses are red* ).
- **Solution:** The negation of  $p$  is given by
- $\sim p$  : *Roses are not red.*
- Therefore. the negation of negation of  $\sim(\sim p)$ :  
*Roses are red.*

# Negation of Implication

- If  $p$  and  $q$  are two statements. Then  $\sim(p \Rightarrow q) \equiv p \wedge \sim q$
- In order to prove the above equivalence. we prepare the following table.

$p$	$q$	$p \Rightarrow q$	$\sim(p \Rightarrow q)$	$\sim q$	$p \wedge \sim q$
T	T	T	F	F	F
T	F	F	T	T	T
F	T	T	F	F	F
F	F	T	F	T	F

- **Example :** Write the negation of each of the following statements:
- If he studies then he will pass in the examination.
- Let  $p : \text{He Studies}$
- $q : \text{He will pass in the examination.}$
- The given statement can be written as  $p \Rightarrow q$ .  
*The negation of  $p \Rightarrow q$  is written as*
- $\sim (p \Rightarrow q) \equiv p \wedge \sim q$ .
- **Result ?????**

# Negation of Biconditional

- If  $p$  and  $q$  are two statements. Then

$$\sim(p \Leftrightarrow q) \equiv p \Leftrightarrow \sim q \equiv \sim p \Leftrightarrow q$$

- **Example** : Write the negation of each of the following statements :
- *He swims if and only if the water is warm.*
- **Solution:** Let  $p$  : *He swims* and  $q$  = *The water is warm.*
- The given statement can be written as  $p \Leftrightarrow q$ .  
*The negation of  $p \Leftrightarrow q$  is written as*

$$\sim(p \Leftrightarrow q) \equiv p \Leftrightarrow \sim q \equiv \sim p \Leftrightarrow q$$

- He swims if and only if the water is not warm.

# Derived Connective

- **NAND** : It means negation of conjunction of two statements. Assume  $p$  and  $q$  be two propositions.
- **NAND** of  $p$  and  $q$  is a proposition which is false when both  $p$  and  $q$  are true otherwise true.
- It is denoted by  $p \uparrow q$ .

- Table: Truth table for NAND

$p$	$q$	$p \uparrow q$
T	T	F
T	F	T
F	T	T
F	F	T

- **NOR:** It means negation of **disjunction** of two statements. Assume  $p$  and  $q$  be two propositions.
- **NOR of  $p$  and  $q$**  is a proposition which is true when both  $p$  and  $q$  are false. Otherwise false. It is denoted by  $p \downarrow q$ .

$p$	$q$	$p \downarrow q$
T	T	F
T	F	F
F	T	F
F	F	T

- **Table: Truth table for NOR**

- **XOR (Exclusive OR):** Assume  $p$  and  $q$  be two proposition. The exclusive or (XOR) of  $p$  and  $q$ , denoted by  $p \oplus q$  is the proposition that is true when exactly one of  $p$  and  $q$  is true but not both and is false otherwise.

- **Table: Truth table for XOR**

$p$	$q$	$p \oplus q$
T	T	F
T	F	T
F	T	T
F	F	F

# Tautologies and Contradictions

- A **tautology** is a proposition which is true for all truth values of its sub-propositions.
- In other words a proposition is a tautology if it is always **true for all assignments** of truth values.
- $\sim(\sim P) \equiv P$

# Contradiction

- A proposition is a **contradiction** if it is always false for all assignments of **truth values**.
- **Remark :**
  - A proposition which is neither a tautology nor a contradiction is called a **contingency**.

# Example

- ***show that  $p \rightarrow p$  is a tautology.***
- **Solution:**

$p$	$p$	$p \rightarrow p$
T	T	T
F	F	T

Since in  $p \rightarrow p$  all the truth values are true (T), hence  $p \rightarrow p$  is a tautology.

# Exercise

1. *Show that  $p \vee (\sim p)$  is a tautology.*
2. *Show that  $\sim(p \wedge (\sim p))$  is a tautology.*
3. *Verify that the proposition  $p \wedge(q \wedge \sim p)$  is a contradiction*

# Functionally Complete Set of Connectives

- Any set of connectives in which every formula can be expressed in terms of an equivalent formula containing the connectives from the set is called a functionally complete set of connectives.
- **Thus** all the conditional and biconditional can be replaced by the three connectives  $\wedge$ ,  $\vee$ ,  $\sim$  .

$$p \Rightarrow q \equiv \sim p \vee q$$

$$p \Leftrightarrow q \equiv (\sim p \vee q) \wedge (p \wedge \sim q)$$

- **Example:** Write an equivalent expression for  $(p \rightarrow q \wedge r) \vee (r \leftrightarrow s)$  which contains neither the biconditional nor the conditional.
- **Solution:** First we replace the biconditional connective by its equivalent in the given expression. Then replace the conditional

$$\begin{aligned}
 (p \Rightarrow q \wedge r) \vee (r \Leftrightarrow s) &\equiv (p \Rightarrow q \wedge r) \vee ((\neg r \vee s) \wedge (r \vee \neg s)) \\
 &\equiv (\neg p \vee q \wedge r) \vee ((\neg r \vee s) \wedge (r \vee \neg s))
 \end{aligned}$$

# Normal Forms

- In logic, with the help of truth table we can compare **if** two statements are equivalent. **But when more statements or propositions are involved, then this method is not practical.**
- One method is to transform  $S_1$  and  $S_2$  to some standard form  $S'_1$  and  $S'_2$ .

# Disjunction Normal Forms

- A logical expression is said to be in disjunctive normal form if it is the sum of elementary products.
- In a logical expression, a **product** of the variables and their negations is called an **elementary product**. For example  $p \wedge \sim q$ ,  $\sim p \wedge \sim q$ ,  $\sim p \wedge q$  are *elementary products*.
- A **sum** of the variables and their negations is called an **elementary sum**. For example,  $p \vee q$ ,  $p \vee \sim q$ ,  $\sim p \vee \sim q$  are *elementary sum*

# Procedure to obtain a disjunctive Normal Form of a given logical expression

1. Remove all  $\rightarrow$  and  $\leftrightarrow$  by an equivalent expression containing the connectives  $\vee, \wedge, \sim$  only.
2. Eliminate  $\sim$  before sums and products by using the double negation.
3. Apply the distributive law until a sum of elementary product is obtained.

- **Example:** Obtain the disjunctive normal forms of the followings:

$$(a) p \wedge (p \Rightarrow q)$$

$$(b) p \vee (\neg p \Rightarrow (q \vee (q \Rightarrow \neg r)))$$

- **Solution:**

a)  $p \wedge (p \Rightarrow q) \equiv [p \wedge (\neg p \vee q) \equiv (p \wedge \neg p) \vee (p \wedge q)]$

b) **homework...**

# Conjunctive Normal Form

- A logical expression is said to be in conjunctive normal form if it consists of a product of elementary sum.
- Example: Obtain a conjunctive normal form of the following:  $[q \vee (p \wedge r)] \wedge \sim [(p \vee r) \wedge q]$
- **Solution:**

$$\begin{aligned}[q \vee (p \wedge r)] \wedge \sim [(p \vee r) \wedge q] &\equiv [q \vee (p \wedge r)] \wedge [\sim (p \vee r) \vee \sim q] \\ &\equiv [q \vee (p \wedge r)] \wedge [(\sim p \wedge \sim r) \vee \sim q] \\ &\equiv (q \vee p) \wedge (q \vee r) \wedge (\sim p \vee \sim q) \wedge (\sim r \vee \sim q)\end{aligned}$$

# Method of Proof

- A **theorem** is a proposition that can be proved to be true.
- An **argument** that establishes the **truth** of a **theorem** is called a **proof**.
- There are **many** different **types** of **proof**.
- In this section we shall look at some of the more common type.

# Rule of Detachment or ( Modus Ponens)

- If the statement in  $p$  is assumed as true and also the statement  $p \rightarrow q$  is accepted as true, then,  $q$  must be true. Symbolically it is written in the following:

$$\frac{p \rightarrow q}{\therefore q}$$

- In this presentation of an argument, the assertions above the horizontal line are the hypotheses or premises; the assertion below the line is the conclusion.

## Example

- If Sushma gets a first class with distinction in B.E. then she will get a good job easily.

Let  $p$  : Sushma gets a first class with distinction in B.E.

$q$  : She will get a good job easily.

$$p \rightarrow q$$

- The inference rule is

$$\frac{p}{\therefore q}$$

Hence this form of argument is valid.

- To do so, we construct a truth table for the premises and conclusion.

$p$	$q$	$p \rightarrow q$	$(p \rightarrow q) \wedge p$	$[(p \rightarrow q) \wedge p] \rightarrow q$
T	T	T	T	T
T	F	F	F	T
F	T	T	F	T
F	F	T	F	T

# Law of Contraposition (or Modus Tollens)

- If the statement in  $p$  is assumed as false and also the statement  $p \rightarrow q$  is accepted as true, then,  $q$  must be false.
- The form of the argument is:  
i.e.  $[(p \rightarrow q) \wedge (\sim q)] \rightarrow \sim p$ .

$$\begin{array}{c} p \rightarrow q \\ \sim q \\ \hline \therefore \sim p \end{array}$$

Truth table for  $[(p \rightarrow q) \wedge (\sim q)] \rightarrow \sim p$ .

$p$	$q$	$\sim p$	$\sim q$	$p \rightarrow q$	$(p \rightarrow q) \wedge (\sim q)$	$[(p \rightarrow q) \wedge (\sim q)] \rightarrow \sim p$
T	T	F	F	T	F	T
T	F	F	T	F	F	T
F	T	T	F	T	F	T
F	F	T	T	T	T	T

# Disjunctive Syllogism

- This rule states that "If  $p \vee q$  is true and  $p$  is false then  $q$  is true."
- It is represented in the following form as

$$\begin{array}{c} p \vee q \\ \hline \neg p \\ \therefore q \end{array}$$

- This argument is valid as  $(p \vee q) \wedge \neg p \rightarrow q$  is a tautology.

p	q	$\neg p$	$(p \vee q)$	$((p \vee q) \wedge \neg p)$	$((p \vee q) \wedge \neg p) \rightarrow q$
T	T	F	T	F	T
T	F	F	T	F	T
F	T	T	T	T	T
F	F	T	F	F	T

# Hypothetical Syllogism

- Whenever the two implications  $p \rightarrow q$  and  $q \rightarrow r$  are accepted as true. then the implication  $p \rightarrow r$  is accepted as true.
- Symbolically it can be represented as

$$\frac{p \Rightarrow q \\ q \Rightarrow r}{\therefore p \Rightarrow r}$$

$p$	$q$	$r$	$p \Rightarrow q$	$q \Rightarrow r$	$p \Rightarrow r$
T	T	T	T	T	T
T	T	F	T	F	F
T	F	T	F	T	T
T	F	F	F	T	F
F	T	T	T	T	T
F	T	F	T	F	T
F	F	T	T	T	T
F	F	F	T	T	T

**Q. Test the validity of the argument. If a person is poor, he is unhappy. If a person is unhappy, he dies young. Therefore poor person dies young.**

Ans. : Let       $p$  : Person is poor

$q$  : Person is unhappy

$r$  : Person dies young

In symbolic form argument is

$$s_1: p \rightarrow q$$

$$\frac{s_2: q \rightarrow r}{s: p \rightarrow r}$$

The above argument is the rule of hypothetical syllogism. Hence it is valid.

# Quiz

**Q. Express following statements in propositional form :**

- i) There are many clouds in the sky but it did not rain.
- ii) I will get first class if and only if I study well and score above 80 in mathematics.
- iii) Computers are cheap but softwares are costly.
- iv) It is very hot and humid or Ramesh is having heart problem.
- v) In small restaurants the food is good and service is poor.
- vi) If I finish my submission before 5.00 in the evening and it is not very hot I will go and play a game of hockey.