



Sorting Algorithms – Merge Sort

Soma Soleiman Zadeh

Data Structures & Algorithms (CBS 216)

Spring 2025 - 2026

Week 13

April 29-30, 2026

Outline



- **Merge Sorting** Algorithm
- **Divide-and-Conquer** Approach
- When to Use **Merge Sort**?
- Implementation of **Merge Sort** in Python
- Time Complexity and Space Complexity

Time Complexity & Space Complexity

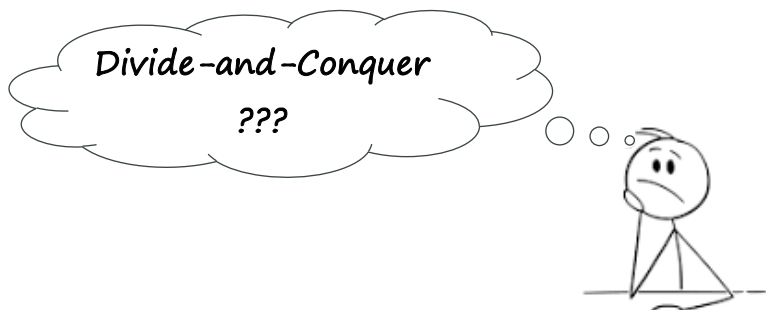


Sorting Algorithm	Time Complexity			Space Complexity
	Best case	Average case	Worst case	Worst case
Selection	$O(n^2)$	$O(n^2)$	$O(n^2)$	$O(1)$
Bubble	$O(n)$	$O(n^2)$	$O(n^2)$	$O(1)$
Insertion	$O(n)$	$O(n^2)$	$O(n^2)$	$O(1)$
Quick	$O(n \log n)$	$O(n \log n)$	$O(n^2)$	$O(n)$
Merge	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	$O(n)$
Heap	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$	$O(1)$

Merge Sort



- The **Merge Sort** is an efficient sorting algorithm that is faster than sorting algorithms like **Bubble sort** and **Insertion sort**.
- The **Merge Sort** follows the **Divide-and-Conquer** approach.

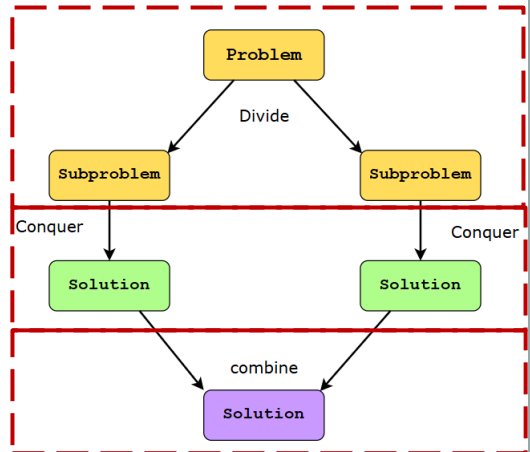




Divide-and-Conquer Approach

◦ The **Divide-and-Conquer** approach is a problem-solving technique that solves a complex problem in three main steps:

1. **Divide:** splitting the main problem into smaller sub-problems.
2. **Conquer:** solving each sub-problem recursively.
3. **Combine:** merging the solutions of sub-problems to form the solution to the original problem.



How Merge Sort Works?

- Merge sort works by recursively dividing the unsorted list into smaller lists until each sub-list contains only one element.
- Then, it merges these sub-lists back into a sorted list. The basic steps are:
 - **Split** the unsorted list into two halves.
 - **Recursively sort each half**.
 - **Merge** the sorted halves into a single sorted list.



How Merge Sort Works?

Unsorted Data

8	2	6	4	5
---	---	---	---	---

Goal



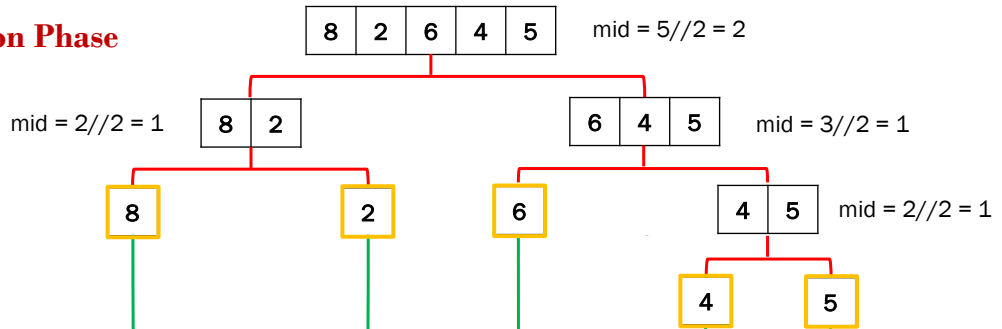
Sorted Data

2	4	5	6	8
---	---	---	---	---

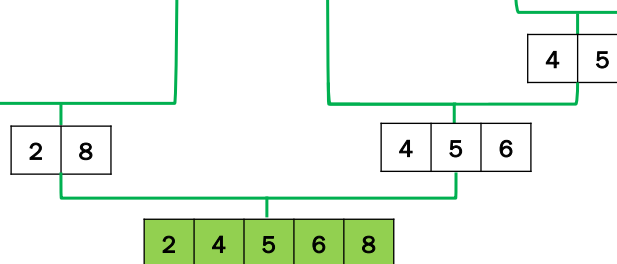
We can show the process of **Merge Sort** algorithm in two phases:

- **Division Phase**
- **Merging Phase**

Division Phase



Merging Phase



When to Use Merge Sort?

- **Merge Sort** is an efficient sorting algorithm where:
 - **Large Datasets** need to be sorted efficiently.
 - **External sorting** is necessary, such as when data is too large to fit into memory (such as RAM).
- **Merge sort** performs far better than **Bubble Sort** and **Insertion Sort**.

Merge Sort Implementation



```
##### Merge Sort Function #####  
  
def mergeSort(data):  
    # Base Case  
    if len(data)<=1:  
        return data  
  
    # Dividing the dataset to Two Halves (Left and Right)  
    mid = len(data)//2  
    left = data[:mid]  
    right = data[mid:]  
  
    # Recursion (Calling the algorithm for each half)  
    mergeSort(left)  
    mergeSort(right)  
  
    # Merging Two Halves  
    i = 0  
    j = 0  
    k = 0  
    while i<len(left) and j<len(right):  
        if left[i]<= right[j]:  
            data[k] = left[i]  
            i += 1  
        else:  
            data[k] = right[j]  
            j += 1  
        k += 1  
  
    # Handling the Remaining Items in Both Halves  
    while i<len(left):  
        data[k] = left[i]  
        i += 1  
        k += 1  
    while j<len(right):  
        data[k] = right[j]  
        j += 1  
        k += 1  
    return data
```

Merge Sort Implementation (Step 1)



```
##### Merge Sort Function #####  
  
def mergeSort(data):  
    # Base Case  
    if len(data)<=1:  
        return data  
  
    # Dividing the dataset to Two Halves (Left and Right)  
    mid = len(data)//2  
    left = data[:mid]  
    right = data[mid:]  
  
    # Recursion (Calling the algorithm for each half)  
    mergeSort(left)  
    mergeSort(right)
```

Merge Sort Implementation (Step 2)



```
# Merging Two Halves  
  
i = 0  
j = 0  
k = 0  
while i<len(left) and j<len(right):  
    if left[i]<= right[j]:  
        data[k] = left[i]  
        i += 1  
    else:  
        data[k] = right[j]  
        j += 1  
  
    k += 1
```

Merge Sort Implementation (Step 3)



```
# Handling the Remaining Items in Both Halves

while i < len(left):
    data[k] = left[i]
    i += 1
    k += 1

while j < len(right):
    data[k] = right[j]
    j += 1
    k += 1

return data
```

Merge Sort: Time and Space Complexity



Sorting Algorithm	Time Complexity (Worst Case)	Space Complexity	Stability
Merge Sort	$O(n \log n)$	$O(n)$	Stable

- **Efficient (Fast) Sorting Algorithm**
- **Good for Sorting Large Datasets**
- **Needs Extra Memory (Out-of-Place Sorting Algorithm)**